



**Smart Contract Audit Report**  
for  
**Almstrong AI**  
Final Report

**July, 2025**

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# 1 Introduction

Trufy has been engaged by what to perform a security audit of the Almstrong AI smart contracts. The purpose of this audit is to achieve the followings:

- Ensure that smart contract functions work as intended.
- Identify possible vulnerabilities, which could be exploited by an attacker.
- Identify smart contract bugs, which might lead to unexpected behavior.
- Make recommendations to improve code safety and readability.

As with any code audit, there is a limit to which vulnerabilities can be found, and unexpected execution paths may still be possible. The author of this report does not guarantee complete coverage.

## 1.1 About Almstrong AI

### 1.1.1 Project Summary

- Project Name: Almstrong AI
- Language: Solidity
- Audit method: Static Analysis, Manual Review
- Scope:
  - ◊ contracts/protocol/adapter
  - ◊ contracts/protocol/lending-vault
  - ◊ contracts/protocol/libraries/logic
  - ◊ contracts/protocol/libraries/tokenization
  - ◊ feature/admin-**contract**/contracts
  - ◊ feature/omni-lending/contracts

## 1.2 Vulnerability Summary

Severity	# of Findings
Critical	3
Medium	1
Low	1
Info	2

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## 2 Findings

ID	Title	Type	Severity	Status
ID-01	Invalid Borrow Validation Using user Instead of onBehalfOf	Logical Issue	Critical	Solved
ID-02	Liquidity Index Overflow via Empty Pool Donation Attack	Logical Issue	Critical	Solved
ID-03	Misuse of LTV as Borrow Ratio Breaks Lending Logic	Logical Issue	Critical	Solved
ID-04	Incorrect Usage of reserve.borrowPool Instead of poolFrom	Logical Issue	Medium	Solved
ID-05	Incorrect Loop Boundary When Accessing Mapping	Logical Issue	Low	Solved
ID-06	Misleading Variable Name onBehalfOf in Collateral Configuration	Informational	Info	Solved
ID-07	Inconsistent Parameter Ordering Between Function and Struct	Informational	Info	Solved

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## 2.1 ID-01: Invalid Borrow Validation Using user Instead of onBehalfOf

Type	Severity	Location	Status
Logical Issue	Critical	CrossChainLendingController.sol#L475	Solved

### 2.1.1 Description

The function `_processValidateBorrowMessage` handles cross-chain borrow requests by decoding message payloads and calling `validateBorrow` to ensure borrowing conditions are met. However, it incorrectly uses `user` instead of `onBehalfOf` when retrieving the borrower's data:

```
440 (vars.user, vars.onBehalfOf, vars.asset, vars.amountToBorrow) =  
    abi.decode(  
441     data,  
442     (address, address, address, uint256)  
443 );  
444  
445 // ...  
446  
447 DataTypes.UserGlobalData storage userData = _users[vars.user];  
448  
449 // ...  
450  
451 ValidationLogic.validateBorrow(  
452     poolData.reserves[vars.asset],  
453     vars.amountToBorrow,  
454     vars.amountInUSD,  
455     _pools,  
456     userData, // @audit this should be userData of `onBehalfOf`  
457     _chainsList,  
458     _chainsCount,  
459     _addressesProvider.getPriceOracle()  
460 );
```

This call uses `userData = _users[vars.user]`, which is the initiator of the message, instead of `_users[vars.onBehalfOf]`, who is the actual target borrower of the requested loan. This creates a **logical inconsistency** where the health factor, collateral, and eligibility checks are performed on the wrong account.

### 2.1.2 Recommendations

- Replace the borrower context in the validation call to use `onBehalfOf`:

---

```
458  DataTypes.UserGlobalData storage userData = _users[vars.  
      onBehalfOf];
```

---

## 2.2 ID-02: Liquidity Index Overflow via Empty Pool Donation Attack

Type	Severity	Location	Status
Logical Issue	Critical	ReserveLogic	Solved

### 2.2.1 Description

When the pool reserve is **empty**, an attacker can exploit the `liquidityIndex` update logic to cause catastrophic overflow in future calculations:

1. Attacker deposits a minimal amount to mint `scaledBalance = 1` unit of TToken.
2. Attacker donates a very large (but tiny in human terms) amount of `aTokens` directly into the pool, e.g.,  $\sim 5e11$  units when the asset uses 18 decimals.
3. The donation drastically increases the `liquidityIndex`, since it is calculated by dividing the donation against the extremely small supply base.

With `totalScaledSupply = 1`, the `liquidityIndex` jumps by  $5e11 * 1e27$ .

4. The system enforces:

```
154 require(newLiquidityIndex <= type(uint128).max, "liquidity  
    index overflow");
```

While the index may not overflow immediately, it can easily surpass safe bounds and overflow **later** when legitimate users have already deposited. This delayed overflow corrupts accounting for **all depositors** and can lock the pool in an inconsistent state.

This creates a **toxic reserve** scenario: the first attacker manipulates the index, and future users are exposed to inevitable accounting failure.

### 2.2.2 Recommendation

Prevent reserves from starting at `scaledSupply = 0`. Specifically:

- Ensure the protocol itself always seeds each reserve with a **non-zero minimum deposit** during initialization.
- Alternatively, enforce a guard that rejects updates when `scaledSupply < threshold` (e.g.,  $< 1e6$  wei).
- Add invariant checks to prevent `liquidityIndex` from increasing by disproportionate amounts relative to real liquidity.

By ensuring the pool is never empty, the exploit vector of artificially inflating the index through small donations is eliminated.

---

## 2.3 ID-03: Misuse of LTV as Borrow Ratio Breaks Lending Logic

Type	Severity	Location	Status
Logical Issue	Critical	BorrowLogic.executeBorrow	Solved

### 2.3.1 Description

In `executeBorrow()`, the function uses `reserve.configuration.getLtv()` from the **borrowed asset** to calculate required collateral:

```
79  uint256 borrowLtv = reserve.configuration.getLtv();
80
81  uint256 collateralNeededInUsd = RefinanceLogic._getUsdValue(
    reserve, params.amount, oracle).percentDiv(borrowLtv);
```

This is incorrect: LTV is defined on **collateral assets**, representing how much can be borrowed *against* them. Using the borrow asset's LTV breaks the lending logic and may allow under-collateralized loans, risking system insolvency.

### 2.3.2 Recommendation

Use the LTV of each collateral asset, not the borrowed asset. Ensure that:

```
1  totalCollateralValue * LTV >= borrowAmount
```

Update the logic accordingly to preserve lending safety.



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## 2.4 ID-04: Incorrect Usage of `reserve.borrowPool` Instead of `poolFrom`

Type	Severity	Location	Status
Logical Issue	Medium	RefinanceLogic.executeRefinanceBorrow	Solved
Logical Issue	Medium	RefinanceLogic.executeRefinanceBorrow	Solved

### 2.4.1 Description

The function `executeRefinanceBorrow` incorrectly uses `reserve.borrowPool` instead of the intended parameter `poolFrom` during the withdrawal process. Specifically, within the `delegatecall` invocation for the `withdraw` method, the code mistakenly references `reserve.borrowPool` instead of the correct `poolFrom`.

This incorrect reference may result in withdrawals from an unintended pool, causing funds to be moved improperly or fail unexpectedly, thereby potentially affecting the refinancing operation.

```
80  uint256 borrowAmount = ILendingAdapter(adapterFrom).
    borrowBalance(reserve.borrowPool, address(this), asset);
    // @audit should be poolFrom
81  if (borrowAmount == 0) return;
82  if (amount > borrowAmount) amount = borrowAmount;
83
84  bool success;
85  for (uint256 i = 0; i < reservesCount; i++) {
86      uint256 maxWithdraw = ILendingAdapter(adapterFrom).
          maxWithdraw(poolFrom, address(this), reserves[i],
            minHf, oracle);
87
88      if (maxWithdraw == 0) continue;
89
90      // withdraw
91      (success,) = adapterFrom.delegatecall(
92          abi.encodeWithSelector(
93              ILendingAdapter.withdraw.selector,
94              reserve.borrowPool, // @audit should be poolFrom
95              reserves[i],
96              maxWithdraw
97          )
98      );
99      require(success, "withdraw failed");
```

---

### 2.4.2 Recommendations

Replace `reserve.borrowPool` with `poolFrom` to ensure withdrawals occur from the correct pool and maintain the intended logical flow of the refinancing operation.

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## 2.5 ID-05: Incorrect Loop Boundary When Accessing Mapping

Type	Severity	Location	Status
Logical Issue	Low	GenericLogic.sol#L209	Solved

### 2.5.1 Description

The function `calculateUserAccountData` iterates over a list of chains stored as a `mapping(uint256 => uint256)` using the loop condition `i <= chainsCount`:

```
209 for (vars.i = 0; vars.i <= chainsCount; vars.i++) {  
210     vars.chainId = chainsList[vars.i];  
211     ...  
212 }
```

In Solidity, accessing a mapping with a key that has never been written to returns the default value `0`. This means that when `vars.i == chainsCount`, the code accesses `chainsList[chainsCount]`, which likely resolves to `0`. Unless `0` is a valid chain ID, this introduces unnecessary computation on default data.

Fortunately, all downstream computations will treat the data associated with chain ID `0` as zeroed values (due to how uninitialized structs behave in Solidity), resulting in no harmful effect on user balances or health factor.

### 2.5.2 Recommendations

- Change the loop condition to use a strict less-than comparison:

```
209 for (vars.i = 0; vars.i < chainsCount; vars.i++) { ... }
```

---

## 2.6 ID-06: Misleading Variable Name onBehalfOf in Collateral Configuration

Type	Severity	Location	Status
Informational	Info	CrossChainLendingController.sol#L135	Solved

### 2.6.1 Description

In the function `_processValidateSetUserUseReserveAsCollateral`, the input data is decoded as follows:

```
135 (address onBehalfOf, address asset, bool useAsCollateral) = abi
    .decode(
136     data,
137     (address, address, bool)
138 );
```

The variable `onBehalfOf` represents the target user whose collateral configuration is being updated. However, the name `onBehalfOf` implies that the action is being performed by a third party on the user's behalf. This is misleading, as there is no delegation or external actor indicated elsewhere in the function — the user in question is the direct subject of the configuration update.

All subsequent interactions refer to this address as the owner of the updated data:

```
147 DataTypes.UserGlobalData storage userData = _users[onBehalfOf];
148 ...
149 UserChainData.userConfig.setUsingAsCollateral(...);
```

### 2.6.2 Recommendations

- Rename the variable `onBehalfOf` to `user` to improve clarity and semantic correctness.
- Avoid naming patterns that suggest proxy or delegated behavior unless such mechanisms are implemented and enforced.

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## 2.7 ID-07: Inconsistent Parameter Ordering Between Function and Struct

Type	Severity	Location	Status
Informational	Info	OmniLendingPool.sol#L206	Solved

### 2.7.1 Description

In the public function `borrow`, the parameters are defined as follows:

```
193 function borrow(  
194     address asset,  
195     uint256 amount,  
196     address onBehalfOf,  
197     uint16 referralCode  
198 )
```

However, when passing these parameters into the `ExecuteBorrowParams` struct for internal processing, the ordering is reversed:

```
210 DataTypes.ExecuteBorrowParams({  
211     asset: asset,  
212     amount: amount,  
213     referralCode: referralCode,  
214     onBehalfOf: onBehalfOf //comes after  
215 });
```

Although the named arguments ensure correct mapping at runtime, this inconsistency introduces cognitive overhead and potential confusion for developers, auditors, and contributors reading the codebase. It also increases the risk of mistakes if the struct is ever instantiated positionally, or in code-generated interfaces and bindings.

### 2.7.2 Recommendations

- Align the order of parameters in the public `borrow()` function to match the struct field order, or vice versa.

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## 3 Appendix

### 3.1 Severity Definitions

#### *Critical*

This level vulnerabilities could be exploited easily and can lead to asset loss, data loss, asset, or data manipulation. They should be fixed right away.

#### *Medium*

This level vulnerabilities are hard to exploit but very important to fix, they carry an elevated risk of smart contract manipulation, which can lead to critical-risk severity.

#### *Low*

This level vulnerabilities should be fixed, as they carry an inherent risk of future exploits, and hacks which may or may not impact the smart contract execution.

#### *Info*

This level vulnerabilities can be ignored. They are code style violations and informational statements in the code. They may not affect the smart contract execution.

### 3.2 Finding Categories

#### ***Gas Optimization***

Gas Optimization findings refer to exhibits that do not affect the functionality of the code but generate different, more optimal EVM opcodes resulting in a reduction on the total gas cost of a transaction.

#### ***Logical Issue***

Logical Issue findings are exhibits that detail a fault in the logic of the linked code, such as an incorrect notion on how block.timestamp works.

#### ***Inconsistency***

Inconsistency findings refer to functions that should seemingly behave similarly yet contain different code, such as a constructor assignment imposing different require statements on the input variables than a setter function.

#### ***Coding Style***

Coding Style findings usually do not affect the generated byte-code and comment on how to make the codebase more legible and as a result easily maintainable.

#### ***Mathematical Operations***

Mathematical Operation exhibits entail findings that relate to mishandling of math formulas, such as overflows, incorrect operations etc.

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### ***Dead Code***

Code that otherwise does not affect the functionality of the codebase and can be safely omitted.

### ***Language Specific***

Language Specific findings are issues that would only arise within Solidity, i.e. incorrect usage of **private** or **delete**.

### ***Centralization***

Centralization findings detail the design choices of designating privileged roles or other centralized controls over the code.